# Binomial Solutions to Smale's 17th Problem and their Application to Chemical Reaction Networks

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July 23, 2014

Key Words

## Table of Contents



2 Binomial Case Solutions

3 Developed Algorithm



Key Words

## Smale's 17th Problem

#### Smale's 17th:

Can one approximate  $\underline{a}$  root of a polynomial system in polynomial-time with a uniform algorithm, on average.

The development of algorithms that are solutions to this problem have applications in real-world polynomial system solving.

#### Previous Work:

- Beltrán & Pardo (2008): Partial positive solution with uniform probabilistic algorithm that runs in time polynomial.
- Qucker & Bürgisser (2012): Deterministic Algorithm running in time N<sup>O(log log N)</sup>.

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### Approximate Root: Smale defined an approximate root as

$$||\zeta - N_f^k(z)|| \le rac{1}{2^{2^k-1}}||\zeta - z||$$

#### Theorem

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Define

$$\gamma(f,\zeta) = \sup_{k \ge 2} \left| \frac{f'(\zeta)^{-1} f^k(\zeta)}{k!} \right|^{\frac{1}{k-1}}$$

Suppose that  $\zeta$  is a true root of f and  $f'(\zeta)^{-1}$  exists. If

$$|z-\zeta|\leq rac{3-\sqrt{7}}{2\gamma(f,\zeta)},$$

then z is an approximate root of f with associated true root  $\zeta$ .

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### Approximate

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### Approximate

- A: The number of distinct roots of a polynomial system can be exponential in the number of variables.
- Olynomial-Time/On Average: Computational complexity takes time polynomial in the input size (number of variables & degree). For a random polynomial, it would run in polynomial time.

### Table of Contents





3 Developed Algorithm



## **Binomial Case Solutions**

#### Lemma

 $x^d-c$  where  $c>0,\ c\in\mathbb{R}_+,\ d\in\mathbb{N},$  and  $d\geq 2$  is a solution to Smale's 17th problem

1 For c > 1, choose interval [0, c] and use bisection method  $\left|x_{n} - |c|^{\frac{1}{d}}\right| < \frac{1}{3d} < \frac{3 - \sqrt{7}}{d - 1}|c|^{\frac{1}{d}}$ 2 For c < 1 choose interval [0, 1] and use bisection method  $\left|x_{n} - |c|^{\frac{1}{d}}\right| < \frac{3 - \sqrt{7}}{d}|c| < \frac{3 - \sqrt{7}}{d - 1}|c|^{\frac{1}{d}}$ 

Complexity:  $O((\log(d)^2))(\log(d) \pm \log(c))$ For a random polynomial, the average value of c is constant.

## Matrix Exponentiation and Smith Normal Form

$$x^{A} := (x_{1}^{A_{11}} ... x_{n}^{A_{n1}}, ..., x_{1}^{A_{1n}} ... x_{n}^{A_{nn}})$$

Then  $x^A = c$  where  $x = (x_1, x_2, \dots x_n)$ ,  $c = (c_1, c_2, \dots c_n)$ , and A is an  $n \times n$  matrix, is written as:

$$x_{1}^{A_{11}}...x_{n}^{A_{n1}} = c_{1}$$
  
:  
$$x_{1}^{A_{1n}}...x_{n}^{A_{nn}} = c_{n}$$

#### The Form

Smith proved that any integer matrix can be diagonalized using elementary row and column operations. SNF will be utilized in the algorithm to solve binomial systems.

## Smith Normal Form

Given

$$\begin{aligned} x_1^2 x_2^4 &= 5\\ x_1^6 x_2^8 &= 9\\ UAV &= S\\ \begin{pmatrix} -1 & 1\\ 2 & -1 \end{pmatrix} \begin{pmatrix} 2 & 6\\ 4 & 8 \end{pmatrix} \begin{pmatrix} 1 & -1\\ 0 & 1 \end{pmatrix} = \begin{pmatrix} 2 & 0\\ 0 & 4 \end{pmatrix}\\ (x^A &= c) \Rightarrow (y^U &= x) \Rightarrow (y^{UA} &= c) \Rightarrow (y^{UAV} &= y^S &= c^V) \end{aligned}$$

Then

$$(y_1)^2 = 5$$
  
 $(y_2)^4 = \frac{9}{5}$ 

 $\begin{array}{l} \mathsf{SNF} \Rightarrow \mathsf{Univariate \ binomials} \Rightarrow \mathsf{Previous \ algorithm} \Rightarrow \\ \tilde{y} \Rightarrow \tilde{y}^U = \tilde{x}. \end{array}$ 

### Table of Contents



2 Binomial Case Solutions

Oeveloped Algorithm



# **Binomial System Algorithm**

### Algorithm

Def binomial system

• Before initiating algorithm, convert binomial system into form  $x^A = c$ 

Input: 
$$A = [a_{ij}] \in \mathbb{Z}^{n \times n}$$
  
Conditions:  $x = [x_1, x_2, \dots, x_n] \ y = [y_1, y_2, \dots, y_n]$   
 $c = (c_1, c_2, \dots, c_n) \in \mathbb{R}^n_+$   
Output: Approximate Root

Calculate the Smith Normal Form of Matrix A. Store
U = [u<sub>ij</sub>] V = [v<sub>ij</sub>] and S = [s<sub>ij</sub>]

Series Perform a monomial change of variable. Let  $y^U = x$ , then  $y^{UA} = c$  and  $y^S = c^V$ 

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# Binomial System Algorithm

### Algorithm

Def binomial system

- Determine  $y^{S}$  and  $c^{V}$  with matrix exponentiation program. Hence,  $y_{1}^{\alpha_{1}} = \beta_{1} \dots y_{n}^{\alpha_{n}} = \beta_{n}$  where  $(\alpha_{1}, \dots, \alpha_{n}) \in \mathbb{N}^{n}$  and  $(\beta_{1}, \dots, \beta_{n}) \in \mathbb{R}^{+}$
- **Solve each equation with bisection method algorithm.** Return vector  $\tilde{y} \in \mathbb{Z}^n$
- Calculate  $\tilde{y}^U$  in order to return variable after monomial change, producing approximate root vector  $\tilde{x}$
- **\bigcirc** Test roots in  $\tilde{x}$  versus approximate root theorem.
- 8 End

## Table of Contents



2 Binomial Case Solutions

3 Developed Algorithm



## Chemical Reaction Networks as Binomial Systems

### Definition

A **Chemical Reaction Network (CRN)** is an organized digraph of chemical reactions. Under the assumption of *mass-action kinetics*, chemical species react at a rate proportional to the product of their concentrations, where the proportionality constant is the rate constant  $\kappa$ .

#### Definition

In chemistry, a steady state is a situation in which all state variables are constant in spite of ongoing processes that strive to change them. It is called a **toric steady state** if the corresponding system of ODE's can be expressed as a binomial system and then solved.

A Chemical Reaction as ODE's

**Reaction** (Triangle Network): Let s = 2 and m = 3. Then  $2A = x_1^2$ ,  $2B = x_2^2$ , and  $A + B = x_1x_2$ 



$$\frac{dx_1}{dt} = -\frac{dx_2}{dt} = (-2\kappa_{12} - \kappa_{13})x_1^2 + (2\kappa_{21} + \kappa_{23})x_2^2 + (\kappa_{31} - \kappa_{32})x_1x_2$$

As only the coefficients of  $x_1x_2$  can be zero, this system has toric steady states iff  $\kappa_{31} = \kappa_{32}$ .

## Chemical Reaction Networks as Binomial Systems

Millán et. al (2011) provided sufficient conditions for a reaction system to have toric steady states.

#### Condition 3.1, 3.4, 3.6

A chemical reaction system is given by a network G with m complexes and reaction rate constants  $\kappa_{ij}^*$ . These three conditions are essentially linear algebra conditions, which result from the concentrations of the chemical species changing over time (the  $\kappa$  values).

## Application Results

Therefore, we can:

- Find biomodels with defined paramaters (κ denotes reaction constants)
- Convert these models' associated dynamical systems into binomial systems based on the work of Millán et. al
- Solution Use binomial system as input for developed binomial system algorithm

The solutions to such systems are of high importance to Computational & Systems Biologists.

References & Acknowledgements

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## Thank You for Listening

#### Acknowledgments:

Kaitlyn Phillipson, *Lead Mentor* Dr. J. Maurice Rojas, *Lead Professor* The National Science Foundation (NSF)

I would like to thank all those who have supported me in my research, the audience members who have taken the time to view today's presentation, and the moderator for their commitment to these presentations.