

Visualizing \mathcal{A} - Discriminants Chambers

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Algorithmic Algebraic Geometry REU

July 23, 2014

Introduction

- We need a software package to visualize the \mathcal{A} -discriminant chambers when $\#\mathcal{A} = n + 4$
- Knowing the chambers helps us easily identify the topological type of the real zero set of real polynomials, among many other applications.

Projective Space

Let \mathbb{K} be any field. We defined projective space \mathbb{P} over the field \mathbb{K} as follows:

$$\mathbb{P}_{\mathbb{K}}^n := \{[z_0 : \cdots : z_n] \mid z_i \in \mathbb{K} \text{ not all zero}\}$$

with the identification

$$[z_0 : \cdots : z_n] = [z_0 \lambda : \cdots : z_n \lambda] \text{ for all } \lambda \in \mathbb{K}^*$$

Horn Kapranov Uniformization

Theorem (M. Kapranov)

Let $\mathcal{A} = \{a_1, \dots, a_{n+k}\} \subset \mathbb{Z}^n$ be the support for a given polynomial f and define the following matrix:

$$\hat{\mathcal{A}} = \begin{pmatrix} 1 & \cdots & 1 \\ a_1 & \cdots & a_{n+k} \end{pmatrix}$$

The parametrization of $\nabla_{\mathcal{A}}$ is given by the closure of

$$\nabla_{\mathcal{A}} = \left\{ [\beta_1 \lambda_1 t^{a_1} : \cdots : \beta_{n+k} \lambda_{n+k} t^{a_{n+k}}] \mid \beta \in \mathbb{C}^{n+k}, \hat{\mathcal{A}}B = 0, t \in (\mathbb{C}^*)^n \right\}$$

where β_i are the rows of B .

Example

$$f(x) = c_0x^7 + c_1x^4 + c_2x^5 + c_3x^3$$

$$\hat{\mathcal{A}} = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 7 & 4 & 5 & 3 \end{pmatrix}$$

$$B = \begin{pmatrix} 1 & 0 \\ 0 & 2 \\ -2 & -1 \\ 1 & -1 \end{pmatrix}$$

$$\nabla_{\mathcal{A}} = \{ [\lambda_1 t^7 : 2\lambda_2 t^4 : (-2\lambda_1 - \lambda_2)t^5 : (\lambda_1 - \lambda_2)t^3] | \lambda_1, \lambda_2 \in \mathbb{C}, t \in \mathbb{C}^* \}$$

Horn Kapranov Uniformization

You can define $\varphi : \mathbb{R}^{k-1} \rightarrow \mathbb{R}^{n+k}$ as $\varphi(\lambda) := \log |\lambda B^T|$

Corollary (Bastani, Hillar, Popov & Rojas)

$$\log |\nabla_{\mathcal{A}}| = \log |\lambda B^T| + \text{row space of } \hat{\mathcal{A}}$$

Example

From...

$$\nabla_{\mathcal{A}} = \{ [\lambda_1 t^7 : 2\lambda_2 t^4 : (-2\lambda_1 - \lambda_2)t^5 : (\lambda_1 - \lambda_2)t^3] | \lambda_1, \lambda_2 \in \mathbb{C}, t \in \mathbb{C}^* \}$$

Taking log

$$\begin{aligned} \text{Log}|\nabla_{\mathcal{A}}| = & \{ (\log |\lambda_1|, \log |2| + \log |\lambda_2|, \log |2\lambda_1 + \lambda_2|, \log |\lambda_1 - \lambda_2|) + \\ & <(1, 1, 1, 1) + (7, 4, 5, 3)> \mid \lambda_1, \lambda_2 \in \mathbb{R} \} \end{aligned}$$

Reduction

Consider $f(x) = c_0 + c_1x + c_2x^2 + c_3x^3$

$$\frac{1}{c_0}f(x) = 1 + \frac{c_1}{c_0}x + \frac{c_2}{c_0}x^2 + \frac{c_3}{c_0}x^3$$

$$\frac{1}{c_0}f\left(\frac{c_0}{c_2}x\right) = 1 + \gamma x + x^2 + \Gamma x^3$$

This can be generalized using the B matrix:

$$c^B = (c_0, c_1, c_2, c_3)^B = \left(\frac{c_0 c_3^2}{c_2^3}, \frac{c_1 c_3}{c_2^2} \right)$$

Reduction

$$\text{Log}|\nabla_{\mathcal{A}}| = \text{Log}|\lambda B^T| + \text{row space of } \hat{\mathcal{A}}$$

We want to consider... $\text{Log}|\bar{\nabla}_{\mathcal{A}}| = \text{Log}|\nabla_{\mathcal{A}}|B$

$$\bar{\varphi} : \mathbb{P}_{\mathbb{R}}^{k-2} \setminus \mathcal{H}_B \rightarrow \mathbb{R}^{k-1} \text{ as } \bar{\varphi}(\lambda) := \log |\lambda B^T|B$$

Reduction

$$f(x) = c_0x^7 + c_1x^4 + c_2x^5 + c_3x^3$$

From...

$$\nabla_{\mathcal{A}} = \{ [\lambda_1 t^7 : 2\lambda_2 t^4 : (-2\lambda_1 - \lambda_2)t^5 : (\lambda_1 - \lambda_2)t^3] | \lambda_1, \lambda_2 \in \mathbb{C}, t \in \mathbb{C}^* \}$$

Taking log

$$\begin{aligned} \text{Log}|\nabla_{\mathcal{A}}| = & \{ (\log |\lambda_1|, \log |2| + \log |\lambda_2|, \log |2\lambda_1 + \lambda_2|, \log |\lambda_1 - \lambda_2|) + \\ & <(1, 1, 1, 1) + (7, 4, 5, 3)> \mid \lambda_1, \lambda_2 \in \mathbb{R} \} \end{aligned}$$

Reduction...

$$\begin{aligned} \text{log } |\bar{\nabla}_{\mathcal{A}}| = & \{ (\log |\lambda_1| + \log |\lambda_1 - \lambda_2| - 2 \log |2\lambda_1 + \lambda_2|, \\ & 2 \log |2| + 2 \log |\lambda_2| - \log |2\lambda_1 + \lambda_2| - \log |\lambda_1 - \lambda_2|) \mid \lambda_1, \lambda_2 \in \mathbb{R} \} \end{aligned}$$

Hyperplane Arrangement

Definition (Hyperplane)

A hyperplane $H \in \mathbb{R}^n$ is any set of the form

$$H = \{x \in \mathbb{R}^n \mid a_1x_1 + \cdots + a_nx_n = c\}$$

for $a \in \mathbb{R}^n$ and real number c .

Definition (Corresponding Hyperplane Arrangement)

$$\mathcal{H}_B = \{[\lambda] \in \mathbb{P}_{\mathbb{R}}^{k-2} \mid \lambda\beta_i = 0 \text{ for some } i \in \{1, \dots, n+k\}\}$$

$$\bar{\varphi}(\lambda) = \log |\lambda B^T| B$$

$$B = \begin{pmatrix} 1 & 0 \\ 0 & 2 \\ -2 & -1 \\ 1 & -1 \end{pmatrix}$$

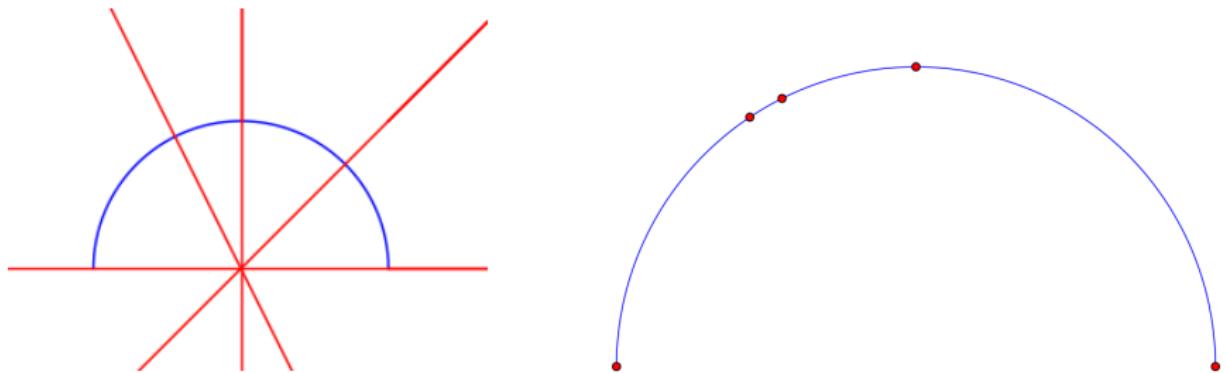
$$\mathcal{H}_B = \{[0 : 1], [1 : 0], [-2 : -1], [1 : -1]\}$$

Lines through the origin in \mathbb{R}^2 where $\bar{\varphi}(\lambda)$ blow up to negative infinity.

n-variate $(n + 3)$ -case: example

Arrangement of points:

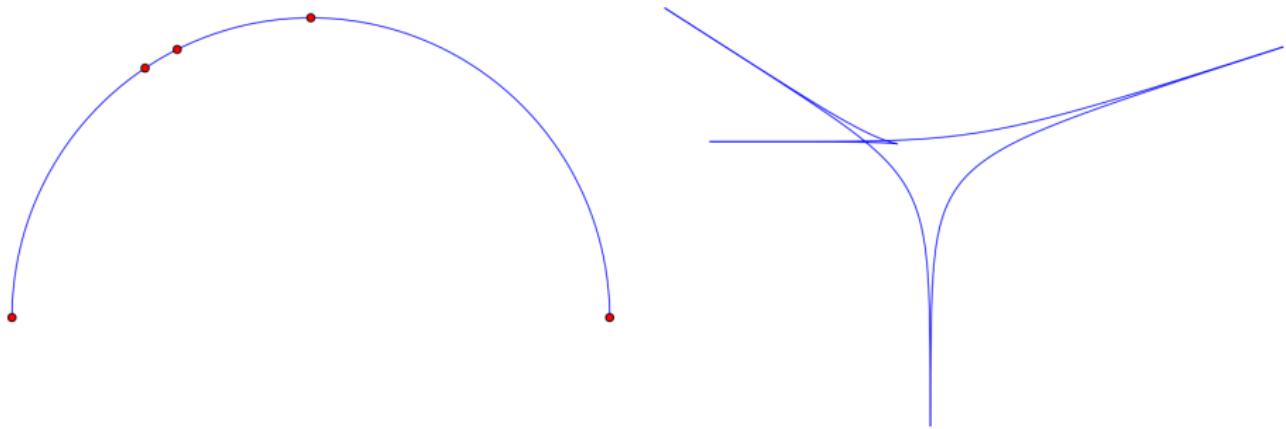
$$f(x) = c_1x^7 + c_2x^4 + c_3x^5 + c_4x^3$$



n-variate $(n + 3)$ -case: example

$$\bar{\varphi}(\lambda) = \log |\lambda B^T |B|:$$

φ maps each arc segment to a piece of the contour amoeba



n-variate $(n+4)$ -case

We are still working on developing a software package for the $(n + 4)$ -case.

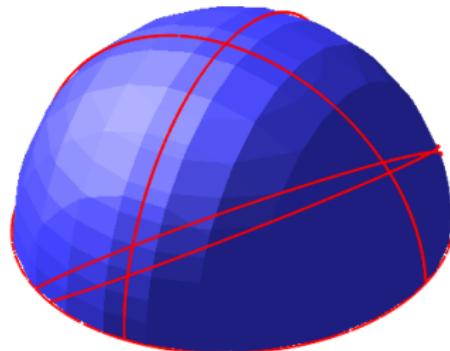
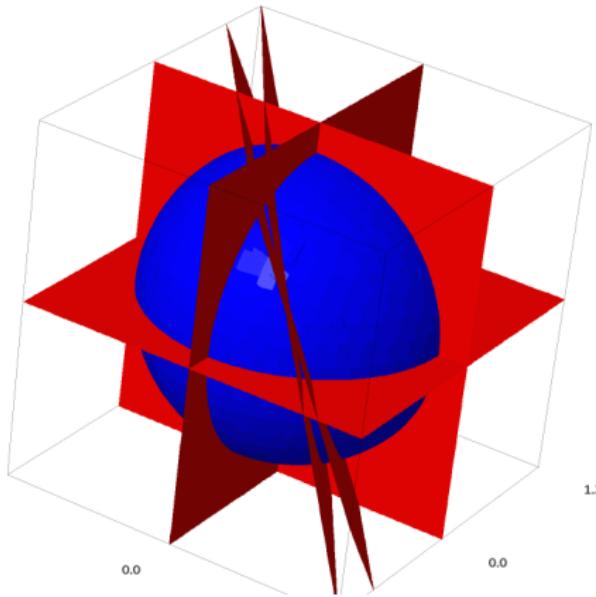
Algorithm (outline)

- Input: Support $\mathcal{A} \subset \mathbb{Z}^n$ with $\#\mathcal{A} = n + 4$
 - Output: Surface of the reduced amoeba in \mathbb{R}^3
- ① Basis for the right null-space of $\hat{\mathcal{A}}$
 - ② Determine the planes where $\bar{\varphi}(\lambda)$ blow up to infinity and represent it as circles (lines) in \mathbb{P}^2
 - ③ Find the intersection lines between the planes and represent it as points in \mathbb{P}^2
 - ④ Store all the information of vertices, edges and faces

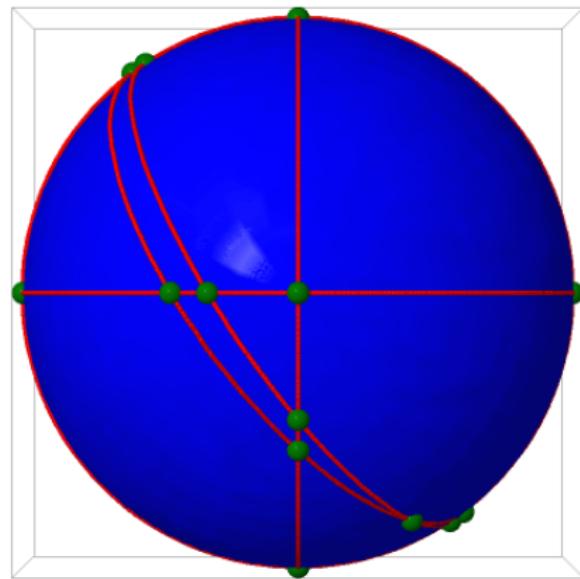
n-variate (n+4)-case: example

Arrangement of lines:

$$f(x) = c_0 + c_1x + c_2x^2 + c_3x^3 + c_4x^4$$



n-variate (n+4)-case: example



References



Kapranov, Misha (1991)

A characterization of A-discriminantal hypersurfaces in terms of the logarithmic Gauss Map

Mathematics Annalen, 290, pp. 277-285



Bastani, Hillar, Popov & Rojas, 2011

Randomization, Sums of Squares, Near-Circuits, and Faster Real Root Counting

Contemporary Mathematics

Thank you for listening!